



MySQL 8 New Features

The Good, The Bad And The Ugly



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Vinicius Grippa is a Percona Senior Support Engineer, Oracle Ace, and author of the book *Learning MySQL*. Vinicius has a Bachelor's degree in Computer Science and has been working with databases for 16 years. He has experience in designing databases for mission-critical applications and, in the last few years, has become a specialist in MySQL and MongoDB ecosystems. Working in the Support team, he has helped Percona customers with hundreds of different cases featuring a vast range of scenarios and complexities. Vinicius is also active in the OS community, participating in virtual rooms like Slack, and speaking at MeetUps, and presenting conferences in Europe, Asia, North and South America.



Agenda

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- ❑ DDL improvements
- ❑ Parallel read threads
- ❑ REDO log capabilities (archiving, disabling)
- ❑ EXPLAIN tree format
- ❑ Options and Variables
- ❑ Performance schema instrumentation
- ❑ Histograms



MySQL 8

MySQL 8

- MySQL 8 Milestone(8.0.0) release 2016-09-12
- MySQL 8 first GA release(8.0.11) is from 2018-04-19
- MySQL is the 2nd in the DB-Engines ranking (behind Oracle)
- Current MySQL GA version 8.0.30 and 8.0.31 being cooked
- Latest Percona Server for MySQL 8.0.29

DDL improvements

ALGORITHM=INSTANT

- Since MySQL 8.0.12, InnoDB supports `ALGORITHM=INSTANT`
- Operations that support `ALGORITHM=INSTANT` only modify metadata in the data dictionary.
- No exclusive metadata locks are taken on the table during preparation and execution phases of the operation, and table data is unaffected, making the operations instantaneous.

ALGORITHM=INSTANT

```
mysql> alter table sbtest2 add g varchar(100) not null default 0,  
algorithm=CONCURRENT;  
Query OK, 0 rows affected (0.50384038 sec)
```

ALGORITHM=INSTANT

- Still requires a metadata lock!

```
1007 | | NULL  
97 | | NULL  
585 | Waiting for table metadata lock | alter table sbtest1 add column test VARCHAR(200) default '', ALGORITHM=INSTANT, LOCK=DEFAULT |  
0 | init | show processlist
```

Parallel threads for online DDL operations

- The number of parallel threads that can be used to scan clustered index is defined by the `innodb_parallel_read_threads` variable. The default setting is 4.
- The number of parallel threads that sort and load data is controlled by the `innodb_ddl_threads` variable, introduced in MySQL 8.0.27. The default setting is 4.
- The new `innodb_ddl_buffer_size` variable defines the maximum buffer size for DDL operations. Defining a buffer size limit avoids potential out of memory errors for online DDL operations that create or rebuild secondary indexes.

Parallel threads for online DDL operations

```
mysql> set session innodb_ddl_threads=16;
mysql> set session innodb_parallel_read_threads=16;
mysql> alter table sbtest2 add g varchar(100) not null default 0,
mysql> algorithm=inplace;
algorithm=inplace;
Query OK, 0 rows affected (1 min 13.90 sec)
Query OK, 0 rows affected (41.66 sec)
```

Parallel read threads

Parallel read threads

- MySQL 8.0.14 has (for now limited) an ability to perform parallel query execution. The `innodb_parallel_read_threads` (default 4) defines the number of threads for parallel index reads.
- It is limited to select `count(*)` from table queries as well as check table queries.

Parallel read threads

```
mysql> set session innodb_parallel_read_threads=1;  
Query OK, 0 rows affected (0.00 sec)
```

```
mysql> select count(*) from joint;
```

```
+-----+
```

```
| count(*) |
```

```
+-----+
```

```
| 134217728 |
```

```
+-----+
```

```
1 row in set (15.29 sec)
```

Parallel read threads

```
mysql> set session innodb_parallel_read_threads=16;  
Query OK, 0 rows affected (0.00 sec)
```

```
mysql> select count(*) from joint;
```

```
+-----+
```

```
| count(*) |
```

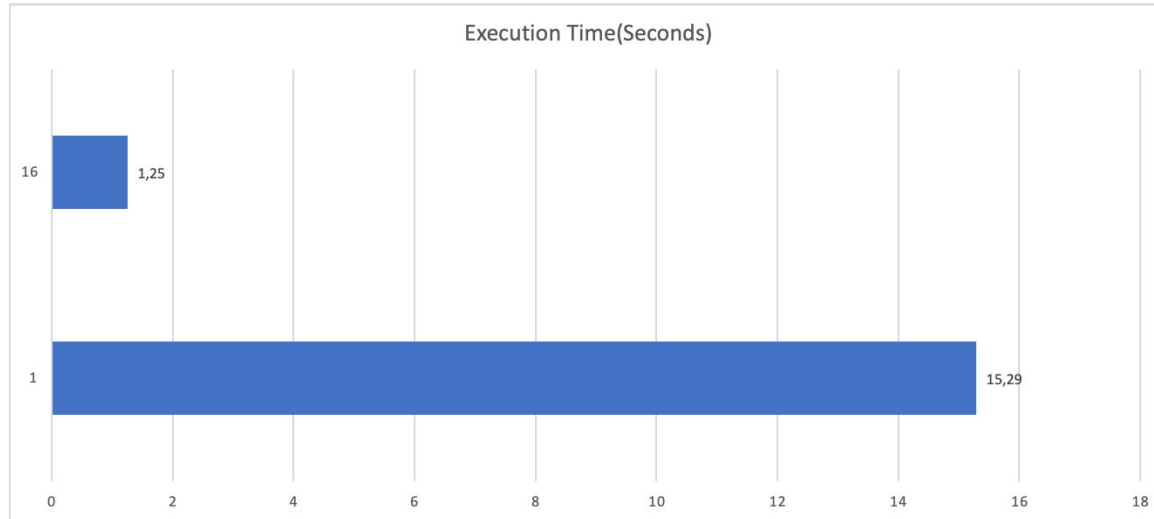
```
+-----+
```

```
| 134217728 |
```

```
+-----+
```

```
1 row in set (1.25 sec)
```

Parallel read threads



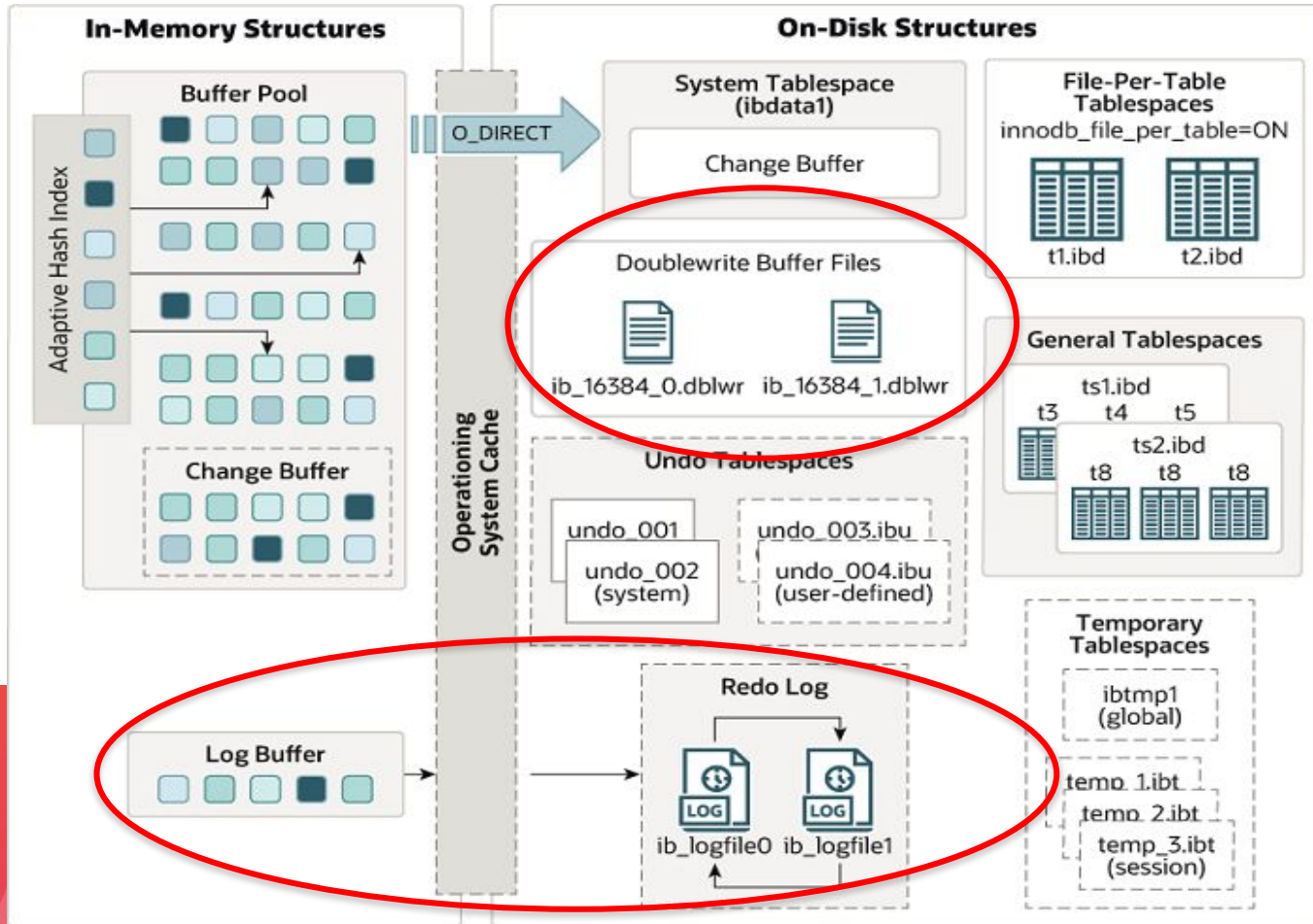
REDO log capabilities

Redo log archiving

- MySQL 8.0.17 releases the redo log archiving(`innodb_redo_log_archive_dir`).
- Backup utilities that copy redo log records may sometimes fail to keep pace with redo log generation while a backup operation is in progress, resulting in lost redo log records due to those records being overwritten.

Disabling redo log

- As of MySQL 8.0.21, we can disable redo logging using the `ALTER INSTANCE DISABLE INNODB REDO_LOG` statement.
- This functionality is intended for loading data into a new MySQL instance. Disabling redo logging speeds up data loading by avoiding redo log writes and doublewrite buffering.



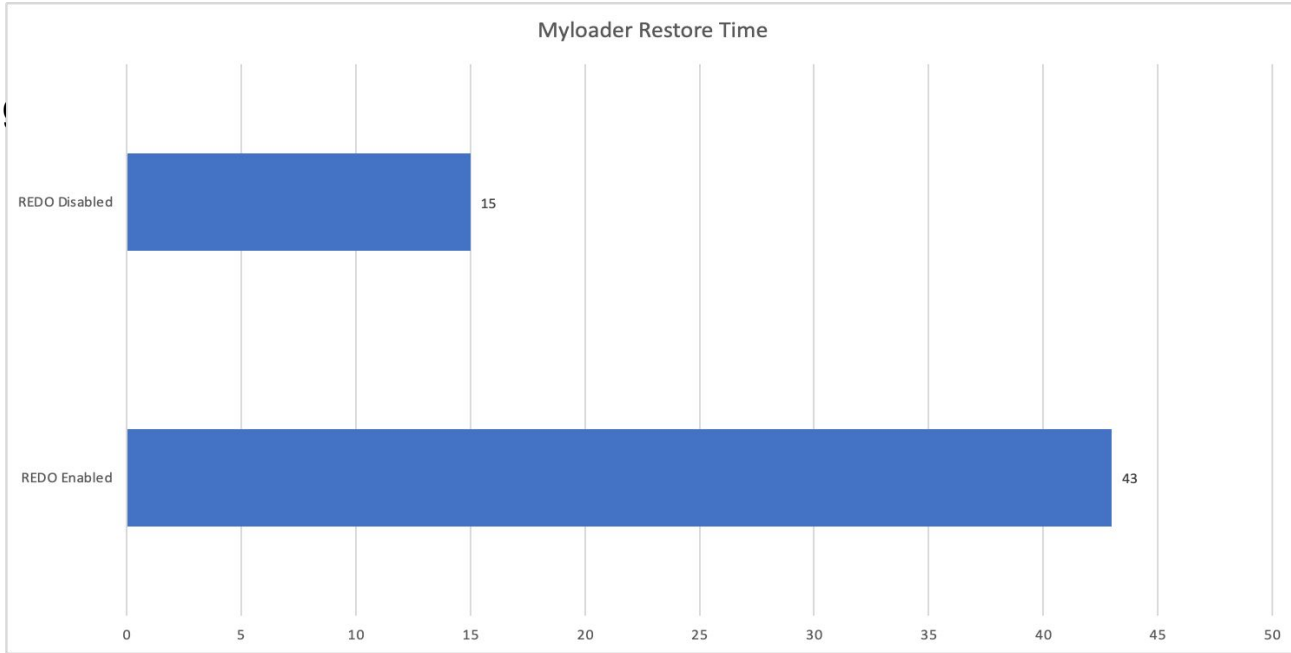
Disabling redo log - benchmark

- **Hardware**
- c5.4xlarge
- 16 CPU
- 32 GB RAM
- 400GB provisioned disk with 3000 IOPs

Disabling redo log - benchmark

- **MySQL 8.0.29**
- 20GB of InnoDB Buffer Pool
- 20GB of redo log size
- sync_binlog=0
- innodb_flush_log_at_trx_commit=0
- ~95GB Data Set

Disabling



Dedicated log writer threads

- As of **MySQL 8.0.22**, you can enable or disable log writer threads using the `innodb_log_writer_threads` variable (default ON).
- Dedicated log writer threads can improve performance on high-concurrency systems, but for low-concurrency systems, disabling dedicated log writer threads provides better performance.

Dedicated log writer threads

- As of **MySQL 8.0.22**, you can enable or disable log writer threads using the `innodb_log_writer_threads` variable (default ON).
- Dedicated log writer threads can improve performance on high-concurrency systems, but for low-concurrency systems, disabling dedicated log writer threads provides better performance.

Dynamic Redo logs

- As of **MySQL 8.0.30**, you can change the redo log capacity dynamically.
- The new `innodb_redo_log_capacity` supersede `innodb_log_file_size` and `innodb_log_files_in_group`.

Dynamic Redo logs

```
2022-08-03T23:03:12.160014Z 0 [Warning] [MY-013869] [InnoDB] Ignored deprecated
configuration parameter innodb_log_file_size. Used innodb_redo_log_capacity
instead.
```

Dynamic Redo logs

```
mysql [localhost:47007] {msandbox} ((none)) > set global
innodb_redo_log_capacity = 2*1024*1024*1024;
Query OK, 0 rows affected (0.23 sec)
```

```
mysql [localhost:47007] {msandbox} ((none)) > show global variables like
'innodb_redo_log_capacity';
```

Variable_name	Value
innodb_redo_log_capacity	2147483648

```
1 row in set (0.01 sec)
```

EXPLAIN tree format

Explain Tree Format

```
EXPLAIN FORMAT=TREE
SELECT first_name, last_name, SUM(amount) AS total
FROM staff INNER JOIN payment
  ON staff.staff_id = payment.staff_id
  AND
  payment_date LIKE '2005-08%'
GROUP BY first_name, last_name;
```

```
-> Table scan on <temporary>
  -> Aggregate using temporary table
    -> Nested loop inner join (cost=1757.30 rows=1787)
      -> Table scan on staff (cost=3.20 rows=2)
        -> Filter: (payment.payment_date like '2005-08%') (cost=117.43 rows=894)
          -> Index lookup on payment using idx_fk_staff_id (staff_id=staff.staff_id) (cost=117.43
rows=8043)
```

Explain tree format

```
EXPLAIN ANALYZE
```

```
SELECT first_name, last_name, SUM(amount) AS total  
FROM staff INNER JOIN payment  
  ON staff.staff_id = payment.staff_id  
  AND  
  payment_date LIKE '2005-08%'  
GROUP BY first_name, last_name;
```

```
-> Table scan on <temporary> (actual time=0.001..0.001 rows=2 loops=1)
```

```
  -> Aggregate using temporary table (actual time=58.104..58.104 rows=2 loops=1)
```

```
    -> Nested loop inner join (cost=1757.30 rows=1787)(actual time=0.816..46.135 rows=5687
```

```
loops=1)
```

```
      -> Table scan on staff (cost=3.20 rows=2)(actual time=0.047..0.051 rows=2 loops=1)
```

```
      -> Filter: (payment.payment_date like '2005-08%') (cost=117.43 rows=894) (actual  
time=0.464..22.767 rows=2844 loops=2)
```

```
        -> Index lookup on payment using idx_fk_staff_id (staff_id=staff.staff_id) (cost=117.43  
rows=8043) (actual time=0.450..19.988 rows=8024 loops=2)
```

Explain tree format

There are several new measurements here:

- Actual time to get first row (in milliseconds)
- Actual time to get all rows (in milliseconds)
- Actual number of rows read
- Actual number of loops

```
Filter: (payment.payment_date like '2005-08%')  
(cost=117.43 rows=894)  
(actual time=0.464..22.767 rows=2844 loops=2)
```

```
1 EXPLAIN ANALYZE  
2 SELECT first_name, last_name, SUM(amount) AS total  
3 FROM staff INNER JOIN payment  
4   ON staff.staff_id = payment.staff_id  
5   AND  
6   payment_date LIKE '2005-08%'  
7 GROUP BY first_name, last_name;  
8  
9 -> Table scan on <temporary> (actual time=0.001..0.001 rows=2 loops=1)  
10   -> Aggregate using temporary table (actual time=58.104..58.104 rows=2 loops=1)  
11   -> Nested loop inner join (cost=1757.30 rows=1787) (actual time=0.816..46.135 rows=5687  
loops=1)  
12   -> Table scan on staff (cost=3.20 rows=2) (actual time=0.047..0.051 rows=2 loops=1)  
13   -> Filter: (payment.payment_date like '2005-08%') (cost=117.43 rows=894) (actual  
time=0.464..22.767 rows=2844 loops=2)  
14   -> Index lookup on payment using idx_fk_staff_id (staff_id=staff.staff_id)  
(cost=117.43 rows=8043) (actual time=0.450..19.988 rows=8024 loops=2)
```

Explain tree format

What can we do with this information?

- If you wonder what it's taking so long, look at the timing. Where does the executor spend its time?
- If you wonder why the optimizer chose that plan, look at the row counters. A large difference (i.e., a couple of orders of magnitude or more) between the estimated number of rows and the actual number of rows is a sign that you should look closer at it. The optimizer chooses its plan based on the estimate, but looking at the actual execution may tell you that another plan would have been better.

Options and variables

innodb_dedicated_server

- Introduced in MySQL 8.0.3

When `innodb_dedicated_server` is enabled, InnoDB automatically configures the following variables:

- `innodb_buffer_pool_size`
- `innodb_log_file_size`
- `innodb_log_files_in_group` (as of MySQL 8.0.14)
- `innodb_flush_method`

innodb_buffer_pool_in_core_file

- Introduced in MySQL 8.0.14
- Disabling the `innodb_buffer_pool_in_core_file` variable reduces the size of core files by excluding InnoDB buffer pool pages.

innodb_use_fdatasync

- Introduced in MySQL 8.0.26
- Allows MySQL to use `ftadasync()` instead of `fsync()`.
- An `fdatasync()` call does not flush changes to file metadata unless required for subsequent data retrieval, providing a potential performance benefit.

innodb_deadlock_detect

- Introduced in MySQL 8.0
- This option is used to disable deadlock detection.

binlog_expire_logs_auto_purge

- Introduced in MySQL 8.0.29
- Enables or disables automatic purging of binary log files (default ON).
- This takes precedence over any setting for `binlog_expire_logs_seconds`

global_connection_memory_limit

- Introduced in MySQL 8.0.28
- The default is virtually infinite (**16 EB**)
- Set the total amount of memory that can be used by all user connections (including `root` user)

global_connection_memory_tracking

- Introduced in MySQL 8.0.28
- Determines whether the server calculates `Global_connection_memory` (Default `FALSE`).

admin_port / admin_address

- Introduced in MySQL 8.0.14
- TCP/IP number to use for connections on administrative interface and IP address to bind for connections.

Performance schema instrumentation

Memory Instrumentation

- Introduced in MySQL 5.7 and enabled by default in MySQL 8.0, the Performance_Schema's Memory instrumentation allows us to have a better overview of what MySQL is allocating and why.

Memory Instrumentation

```
SELECT SUBSTRING_INDEX(event_name, '/', 2) AS code_area,  
       sys.format_bytes(SUM(current_alloc)) AS current_alloc  
FROM sys.x$memory_global_by_current_bytes  
GROUP BY SUBSTRING_INDEX(event_name, '/', 2)  
ORDER BY SUM(current_alloc) DESC;
```

Memory Instrumentation

```
+-----+-----+
| code_area          | current_alloc |
+-----+-----+
| memory/innodb      | 21.09 GiB     |
| memory/performance_schema | 236.14 MiB    |
| memory/sql         | 36.26 MiB     |
| memory/mysys       | 9.43 MiB      |
| memory/temptable  | 1.00 MiB      |
| memory/mysqld_openssl | 604.85 KiB    |
| memory/csv         | 18.07 KiB     |
| memory/mysqllx     | 3.31 KiB      |
| memory/myisam      | 728 bytes     |
| memory/blackhole   | 120 bytes     |
| memory/vio         | 80 bytes      |
+-----+-----+
11 rows in set, 1 warning (0.03 sec)
```

TempTable Memory Usage

```
mysql>select * from sys.memory_global_by_current_bytes where event_name like  
'%temp%'\G
```

Memory Usage Per User

```
mysql> select * from sys.memory_by_user_by_current_bytes\G
```

Memory Usage Per User

```
***** 1. row *****
      user: root
current_count_used: 13094
  current_allocated: 4.51 MiB
  current_avg_alloc: 361 bytes
  current_max_alloc: 1.29 MiB
    total_allocated: 15.33 TiB
***** 2. row *****
      user: sysbench
current_count_used: 0
  current_allocated: 0 bytes
  current_avg_alloc: 0 bytes
  current_max_alloc: 0 bytes
    total_allocated: 4.34 TiB
2 rows in set (0.00 sec)
```

Memory Usage Per Thread

```
mysql> select * from memory_by_thread_by_current_bytes\G
```

Memory Usage Per Thread

```
***** 1. row *****
      thread_id: 1072
        user: root@localhost
current_count_used: 13100
  current_allocated: 11.57 MiB
  current_avg_alloc:  925 bytes
  current_max_alloc: 8.00 MiB
    total_allocated: 152.40 MiB
***** 2. row *****
      thread_id: 1
        user: sql/main
current_count_used: 6552
  current_allocated: 1.14 MiB
  current_avg_alloc:  182 bytes
  current_max_alloc: 357.13 KiB
    total_allocated: 6.04 MiB
```

Histograms

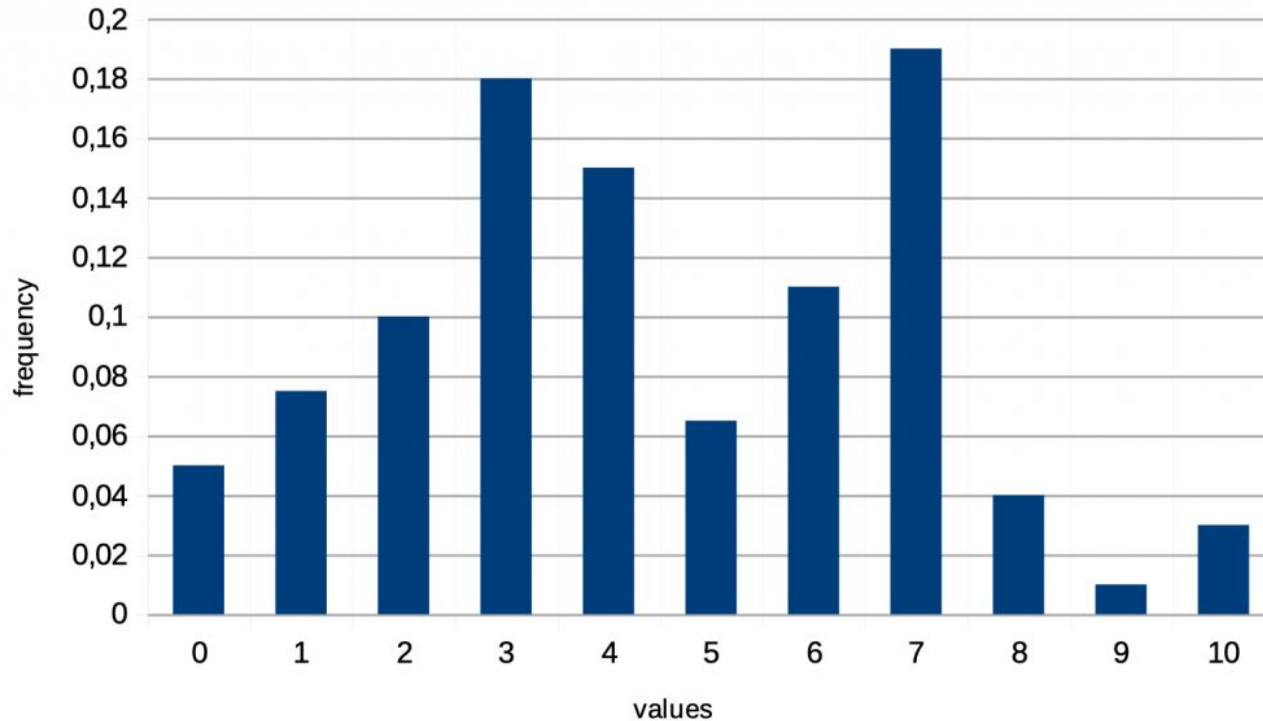
Histograms

- We can define a histogram as a good approximation of the data distribution of the values in a column.
- So in MySQL, histograms help the optimizer to find the most efficient query Plan.

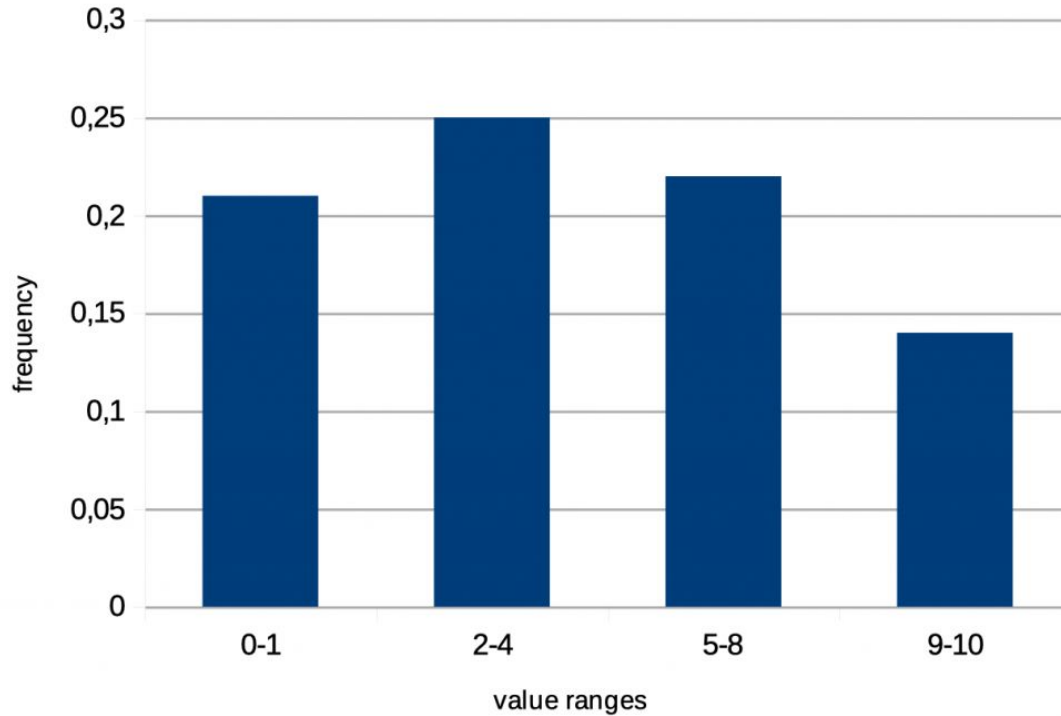
Histograms

- Histograms can be of two types
 - Singleton
 - Equi-height

Histograms - Singleton



Histograms - Equi-height



Histograms

```
mysql> ANALYZE TABLE country UPDATE HISTOGRAM ON Region;
```

Histograms

```
mysql> > SELECT SUBSTRING_INDEX(v, ':', -1) value, concat(round(c*100,1),'%') cumulfreq,  
CONCAT(round((c - LAG(c, 1, 0) over()) * 100,1), '%') freq FROM information_schema.column_statistics,  
JSON_TABLE(histogram->>'$.buckets','$[*]' COLUMNS(v VARCHAR(60) PATH '$[0]', c double PATH '$[1]')) hist  
WHERE schema_name = 'world' and table_name = 'country' and column_name = 'region';
```

value	cumulfreq	freq
Antarctica	2.1%	2.1%
Australia and New Zealand	4.2%	2.1%
[...]		
Central Africa	20.1%	3.8%
Central America	23.4%	3.3%
Eastern Africa	31.8%	8.4%
[...]		
Western Africa	96.2%	7.1%
Western Europe	100%	3.8%

25 rows in set (0.00 sec)

Histograms

- Histogram statistics are particularly useful for non-indexed columns, as shown in the example.
- Execution plans that can rely on indexes are usually the best, but histograms can help in some edge cases or when creating a new index is a bad idea.

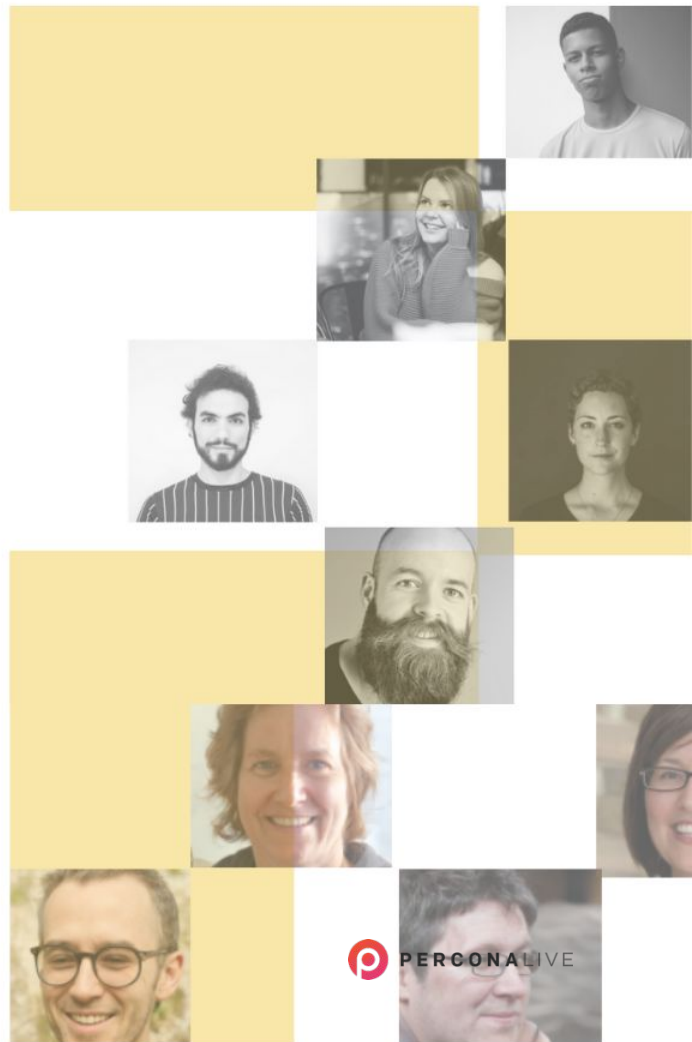
We're hiring
Join our team!

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APPLY NOW: percona.com/careers



谢谢

Thank you

Grazie

Obrigado

Gracias

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